

Corrections to  
*Set Theory and the Continuum Problem* (revised edition)

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**These are corrections to the edition published by Dover in 2010.**

**Page 23** (Error found by David Feuer) Exercise 5.5(d) in Chapter 2 asks the reader to show  $B - (A - B) = \emptyset$ . It should be to show  $B - (A - B) = B$ .

**Page 23** (Error found by David Feuer) Text in §6 reads  $A_6$  [**Power set axiom**]. It should read  $A_6$  [**Power set axiom**].

**Page 25** (Error found by Gavin Rebeiro) On pg. 25 line 36 (5th line of last paragraph), “if  $x \neq y$  then  $F(x) \neq F(Y)$ ” should be “...  $F(x) \neq F(y)$ ”.

**Page 39** (Error found by Gavin Rebeiro) On pg. 36, last paragraph on page. “Let  $L$  be the class of all elements  $x$  such that  $x$  is a proper subclass of all elements of  $A$ ” should be “Let  $L$  be the class of all elements  $x$  in  $M$  such that  $x$  is a proper subclass of all elements of  $A$ .”

**Page 51** (Error found by Grigory Olkhovikov) Exercise 1.2 is incorrect as asked. Take  $A$  to be the set of negative integers under the natural linear order  $\leq$ . Every proper lower section of  $A$  has a strict upper bound, but its ordering is not a well ordering. The exercise can be corrected by including the condition that  $A$  has a least element

**Page 55** (Error found by Hal Prince) line 6: “every  $y$  in  $M$ ” should be “every  $y$  in  $C$ ”.

**Page 60** (Error found by Allen David Boozer) Exercise 4.3 should be deleted. There is a reference to this Exercise in the Remark at the top of page 63, and this reference should also be deleted.

**Page 66** (Error found by Stuart Newberger) Definition 7.1 is incorrect as stated. It should read as follows.

For any sets  $y$  and  $x$ , we will say that  $y$  is *closed* (under  $g$ ) *relative to*  $x$  provided, for any  $z \in y$ , if  $g(z) \in \mathcal{P}(x)$  then  $g(z) \in y$ . (Thus  $(z \in y \wedge g(z) \subseteq x) \supset g(z) \in y$ .)

**Page 225-226** (Error found by Chang Soon Choi.) In the Remarks, it is not the case that  $(f \approx_\lambda g) \supset \llbracket f \approx_\lambda g \rrbracket$  is **S4** valid generally, but it is valid in the particular **S4** models being constructed. Here is the argument. Suppose as an induction hypothesis that it is known for ordinals less than  $\lambda$ . It follows from the definition of  $\approx_\lambda$  that if  $\alpha < \lambda$  then  $(f \approx_\alpha g) \supset (f \approx_\lambda g)$  is valid in these models. It follows from this, using general **S4** reasoning, that  $\Box\Diamond(f \approx_\alpha g) \supset \Box\Diamond(f \approx_\lambda g)$  is also valid in these models, that is,  $\llbracket f \approx_\alpha g \rrbracket \supset \llbracket f \approx_\lambda g \rrbracket$ . Now if  $p \Vdash (f \approx_\lambda g)$ , then  $p \Vdash (f \approx_\alpha g)$  for some  $\alpha < \lambda$  (by definition). This implies  $p \Vdash \llbracket f \approx_\alpha g \rrbracket$ , and hence  $p \Vdash \llbracket f \approx_\lambda g \rrbracket$ .

**Page 227** (Error found by Chang Soon Choi.) In the proof of Proposition 1.5, the limit ordinal case is incorrect. It uses an inference from  $p \Vdash \llbracket f \approx_\lambda g \rrbracket$ , where  $\lambda$  is a limit ordinal, to  $p \Vdash \llbracket f \approx_\alpha g \rrbracket$ , for some  $\alpha < \lambda$ , and this is not justified. Replace the limit ordinal case by the following.

Assume  $\lambda$  is a limit ordinal and every  $\alpha < \lambda$  is good. Now suppose  $p \Vdash \llbracket f \approx_\lambda g \rrbracket$  and  $\lambda < \beta$ ; we must show  $p \Vdash \llbracket f \approx_\beta g \rrbracket$ .

We first show  $(f \approx_\lambda g) \supset \llbracket f \approx_\beta g \rrbracket$  is valid in the model that has been constructed (note the absence of double square brackets in the antecedent). Well, suppose  $q \Vdash (f \approx_\lambda g)$ . Then  $q \Vdash (f \approx_\alpha g)$  for some  $\alpha < \lambda$ , by definition of  $\approx_\lambda$ . It follows by the Remarks at the bottom of page 225 and the top of 226 that  $q \Vdash \llbracket f \approx_\alpha g \rrbracket$ . Since  $\alpha$  is good,  $q \Vdash \llbracket f \approx_\beta g \rrbracket$ . Since  $q$  was arbitrary, we have shown the validity of  $(f \approx_\lambda g) \supset \llbracket f \approx_\beta g \rrbracket$  in the model.

It now follows, by standard modal manipulations, that  $\Box\Diamond(f \approx_\lambda g) \supset \Box\Diamond\llbracket f \approx_\beta g \rrbracket$  is also valid in the model, and hence we have the validity of  $\llbracket f \approx_\lambda g \rrbracket \supset \llbracket f \approx_\beta g \rrbracket$ , making use of Proposition 4.3, part 2. Since  $p \Vdash \llbracket f \approx_\lambda g \rrbracket$ , then  $p \Vdash \llbracket f \approx_\beta g \rrbracket$ .

**Page 228** (Found by Grigori Mints) In the Remark just before Definition 1.8 it is asserted that  $(f \in g) \equiv \llbracket f \in g \rrbracket$ . The equivalence is not correct, but  $(f \in g) \supset \llbracket f \in g \rrbracket$  is.

**Page 229** (Problem found by Chang Soon Choi.) Lemma 2.1 says that if  $p \Vdash \llbracket f \approx_\alpha g \rrbracket$  and  $p \Vdash \llbracket g \approx_\alpha h \rrbracket$  then  $p \Vdash \llbracket f \approx_\alpha h \rrbracket$ . The proof is by induction on  $\alpha$ . It begins by saying the cases where  $\alpha$  is 0 or a limit ordinal are simple. In fact 0 is simple, but the limit ordinal case is not. Here is a proof for the limit ordinal case.

Let  $\lambda$  be a limit ordinal and assume the result holds for smaller ordinals. We begin by showing that if  $p \Vdash (f \approx_\lambda g)$  and  $p \Vdash \llbracket g \approx_\lambda h \rrbracket$  then  $p \Vdash \llbracket f \approx_\lambda h \rrbracket$  (note the difference in the first item). So, suppose  $p \Vdash (f \approx_\lambda g)$  and  $p \Vdash \llbracket g \approx_\lambda h \rrbracket$ . To show  $p \Vdash \llbracket f \approx_\lambda h \rrbracket$  we show  $p \Vdash \Box\Diamond(f \approx_\lambda h)$ . Let  $q$  be any member of  $\mathcal{G}$  such that  $p\mathcal{R}q$ ; we must show there is some  $r$  with  $q\mathcal{R}r$  so that  $r \Vdash (f \approx_\lambda h)$ .

Since  $p \Vdash \Box\Diamond(g \approx_\lambda h)$  then  $q \Vdash \Diamond(g \approx_\lambda h)$  and hence there is some  $r$  with  $q\mathcal{R}r$  so that  $r \Vdash (g \approx_\lambda h)$ . By definition,  $r \Vdash (g \approx_\alpha h)$  for some  $\alpha < \lambda$ . Without loss of generality we can assume  $\alpha$  is a successor ordinal. Then  $r \Vdash \llbracket g \approx_\alpha h \rrbracket$  by the remarks on pages 225-226.

Since  $p \Vdash (f \approx_\lambda g)$  then  $p \Vdash (f \approx_\beta g)$  for some  $\beta < \lambda$  and again without loss of generality we can assume  $\beta$  is a successor ordinal. Then  $p \Vdash \llbracket f \approx_\beta g \rrbracket$  by the remarks on pages 225-226 again. Since this formula begins with  $\Box$ ,  $r \Vdash \llbracket f \approx_\beta g \rrbracket$ . Let  $\gamma$  be the larger of  $\alpha$  and  $\beta$ . By Proposition 1.5,  $r \Vdash \llbracket f \approx_\gamma g \rrbracket$  and  $r \Vdash \llbracket g \approx_\gamma h \rrbracket$ . Since  $\gamma < \lambda$ , by the induction hypothesis for the overall Lemma,  $r \Vdash \llbracket f \approx_\gamma h \rrbracket$ . Since  $\gamma$  is a successor ordinal, by the remarks on pages 225-226 again,  $r \Vdash (f \approx_\gamma h)$ , which is what we wanted.

Since  $p$  was arbitrary, we have shown the validity in our model of

$$(f \approx_\lambda g) \supset (\llbracket g \approx_\lambda h \rrbracket \supset \llbracket f \approx_\lambda h \rrbracket).$$

Then by standard S4 manipulations, this gives us the validity in our model of

$$\Box\Diamond(f \approx_\lambda g) \supset \Box\Diamond(\llbracket g \approx_\lambda h \rrbracket \supset \llbracket f \approx_\lambda h \rrbracket).$$

By Proposition 4.4 of Chapter 16 we then have

$$\Box\Diamond(f \approx_\lambda g) \supset \llbracket g \approx_\lambda h \supset f \approx_\lambda h \rrbracket$$

and hence

$$\llbracket f \approx_\lambda g \rrbracket \supset (\llbracket g \approx_\lambda h \rrbracket \supset \llbracket f \approx_\lambda h \rrbracket)$$

by using Proposition 4.5 of Chapter 16.

**Page 234** (Problem found by Chang Soon Choi.) In line 8 of the proof of Proposition 3.3, “equivalently,  $\llbracket \hat{s} \in \hat{t} \rrbracket$ ” should be changed to “and so  $\llbracket \hat{s} \in \hat{t} \rrbracket$ ”. Also in line 11 of the same proof, “But then  $p \Vdash (a \varepsilon \hat{t})$ , so  $a$  is  $\hat{x} \dots$ ” should be changed to “So  $a$  is  $\hat{x} \dots$ ”.

**Pages 239–240** (Problem found by Chang Soon Choi.) In the proof of Lemma 5.2 it is said that “(Recall that  $(x \approx_\alpha y)$  and  $\llbracket x \approx_\alpha y \rrbracket$  are equivalent.)” This is not the case. One should modify the condition that we need to express by a first-order formula so that the last part reads  $\Box \Diamond (x \approx_\alpha y)$ . Then, in the formula following “Now let  $F(\mathcal{A}, p, f, g)$  be the formula:” the final clause should be changed from “ $\langle s', x, y \rangle \in \mathcal{A}$ ” to “ $\langle r', x, y \rangle \in \mathcal{A}$ ”.

**Page 241** (Problem found by Chang Soon Choi.) In the proof of Theorem 5.4, the atomic case should be modified. We know that  $p \Vdash \llbracket f \in g \rrbracket$  iff  $p \Vdash \llbracket (\exists w)(w \approx f \wedge w \varepsilon g) \rrbracket$  iff  $p \Vdash \Box \Diamond (\exists w)(\llbracket w \approx f \rrbracket \wedge \llbracket w \varepsilon g \rrbracket)$ , so in the atomic case,  $F_\varphi(z, x, y)$  should be

$$(\forall z' \overleftarrow{R} z)(\exists z'' \overleftarrow{R} z')(\exists w \in \mathcal{D})(\text{Equals}(z'', w, x) \wedge (\forall z''' \overleftarrow{R} z'')(\exists z'''' \overleftarrow{R} z''')(\langle z''''', w \rangle \in y))$$

Also in the final part of the proof take  $F_\varphi(z, x_1, \dots, x_n)$  to be the following:

$$(\forall z' \overleftarrow{R} z)(\exists z'' \overleftarrow{R} z') \neg F_\psi(z'', x_1, \dots, x_n).$$

**Page 264** (Problem found by Jason Parker) The remarks at the end of the first paragraph are incorrect. First, a few useful observations: using Definition 3.1 on page 233, one has the following.

$$\begin{aligned} \hat{0} &= \emptyset \\ \hat{1} &= \mathcal{G} \times \{\hat{0}\} \\ \hat{2} &= \mathcal{G} \times \{\hat{0}, \hat{1}\} \\ &= \hat{1} \cup (\mathcal{G} \times \{\hat{1}\}) \\ \hat{3} &= \mathcal{G} \times \{\hat{0}, \hat{1}, \hat{2}\} \\ &= \hat{2} \cup (\mathcal{G} \times \{\hat{2}\}) \\ &\vdots \\ \hat{\omega} &= \mathcal{G} \times \{\hat{0}, \hat{1}, \hat{2}, \dots\} \end{aligned}$$

Now, here is Parker’s argument.

“It is claimed that we can show that if  $p \Vdash \llbracket a \subseteq \hat{\omega} \rrbracket$ , then for some  $b \subseteq \mathcal{G} \times \hat{\omega}$ ,  $p \Vdash \llbracket a \approx b \rrbracket$ . But this would entail that  $b \in \mathcal{D}^\mathcal{G}$ , which does not seem possible. For suppose  $b \subseteq \mathcal{G} \times \hat{\omega}$ . If  $b$  were a member of  $\mathcal{D}^\mathcal{G}$ , then  $b \in R_{\alpha+1}^\mathcal{G}$  for some least ordinal  $\alpha$ . Then  $b \subseteq \mathcal{G} \times R_\alpha^\mathcal{G}$ . Now since  $b \subseteq \mathcal{G} \times \hat{\omega}$ , it follows that any  $x \in b$  is of the form  $\langle p, \langle q, \hat{n} \rangle \rangle$  for some  $p, q \in \mathcal{G}$  and  $n \in \omega$ . So  $\langle q, \hat{n} \rangle \in R_\alpha^\mathcal{G}$ . So there is some least ordinal  $\beta$  such that  $\langle q, \hat{n} \rangle \in R_{\beta+1}^\mathcal{G}$ , whereby  $\langle q, \hat{n} \rangle \subseteq \mathcal{G} \times R_\beta^\mathcal{G}$ , which is clearly false. So it seems that it cannot be that  $b \in \mathcal{D}^\mathcal{G}$  if  $b \subseteq \mathcal{G} \times \hat{\omega}$ .”

The problem sentences at the end of paragraph 1, page 264, should be replaced with the following. “Now, this result can be improved, to establish that if  $\llbracket a \subseteq \hat{\omega} \rrbracket$  is true at  $p$  then

$\llbracket a \approx b \rrbracket$  is true at  $p$  for some  $b \subseteq \mathcal{G} \times \{\hat{n} \mid n \in \omega\}$  (equivalently, for some  $b \subseteq \hat{\omega}$ ). Consequently, to investigate the size of the power set of  $\hat{\omega}$  in the modal model, we begin by investigating the actual power set of  $\hat{\omega}$ .”

Here is the argument for the revised assertion above. Throughout, assume that  $p \Vdash \llbracket a \subseteq \hat{\omega} \rrbracket$ , meaning  $p \Vdash \llbracket (\forall x)(x \in a \supset x \in \hat{\omega}) \rrbracket$ .

1. If  $p\mathcal{R}p'$  and  $p' \Vdash \llbracket x \in a \rrbracket$ , then for some  $p''$  with  $p'\mathcal{R}p''$ ,  $p'' \Vdash \llbracket x \approx \hat{n} \wedge \hat{n} \in a \rrbracket$  for some  $n \in \omega$ .  
 Proof: Suppose  $p\mathcal{R}p'$ , and  $p' \Vdash \llbracket x \in a \rrbracket$ . Then  $p' \Vdash \llbracket x \in \hat{\omega} \rrbracket$ , and so for some  $p''$  with  $p'\mathcal{R}p''$ ,  $p'' \Vdash x \in \hat{\omega}$ , and hence for some  $h$ ,  $p'' \Vdash \llbracket x \approx h \rrbracket$  and  $p'' \Vdash \llbracket h \varepsilon \hat{\omega} \rrbracket$  (Definition 1.6 Chapter 17). Then for some  $p'''$  with  $p''\mathcal{R}p'''$ ,  $p''' \Vdash h \varepsilon \hat{\omega}$ , and so  $\langle p''', h \rangle \in \hat{\omega} = \mathcal{G} \times \{\hat{0}, \hat{1}, \dots\}$ . Then  $h = \hat{n}$  for some  $n \in \omega$ . It follows that  $p'' \Vdash \llbracket x \approx \hat{n} \rrbracket$  and  $p'' \Vdash \llbracket \hat{n} \in a \rrbracket$ .
2. Now let  $b = \{\langle q, \hat{n} \rangle \mid n \in \omega, q \Vdash \llbracket \hat{n} \in a \rrbracket\}$ . Trivially  $b \subseteq \mathcal{G} \times \{\hat{0}, \hat{1}, \dots\} = \hat{\omega}$ .
3.  $p \Vdash \llbracket a \subseteq b \rrbracket$ . The proof is by contradiction. Suppose not; then for some  $h$  and for some  $p'$  with  $p\mathcal{R}p'$ ,  $p' \Vdash \llbracket h \in a \rrbracket$  and  $p' \Vdash \llbracket \neg(h \in b) \rrbracket$  (**P**<sub>9</sub>, Page 226). By item 1, for some  $p''$  with  $p'\mathcal{R}p''$ ,  $p'' \Vdash \llbracket h \approx \hat{n} \rrbracket$  and  $p'' \Vdash \llbracket \hat{n} \in a \rrbracket$  for some  $n \in \omega$ . Let  $q$  be an arbitrary member of  $\mathcal{G}$  with  $p''\mathcal{R}q$ . Then  $q \Vdash \llbracket \hat{n} \in a \rrbracket$ , hence  $\langle q, \hat{n} \rangle \in b$ , and so  $q \Vdash \hat{n} \varepsilon b$ . Since  $q$  was arbitrary,  $p'' \Vdash \Box(\hat{n} \varepsilon b)$ , and so  $p'' \Vdash \Box\Diamond(\hat{n} \varepsilon b)$ , or  $p'' \Vdash \llbracket \hat{n} \varepsilon b \rrbracket$ . Then  $p'' \Vdash \llbracket \hat{n} \in b \rrbracket$  (Corollary 2.4, Chapter 17). But we also have  $p'' \Vdash \llbracket \neg(\hat{n} \in b) \rrbracket$ , and this is our contradiction.
4.  $p \Vdash \llbracket b \subseteq a \rrbracket$ . Again the proof is by contradiction. If not, then for some  $h$  and for some  $p'$  with  $p\mathcal{R}p'$ ,  $p' \Vdash \llbracket h \in b \rrbracket$  and  $p' \Vdash \llbracket \neg(h \in a) \rrbracket$ . Then for some  $p''$  with  $p'\mathcal{R}p''$ ,  $p'' \Vdash h \in b$  and hence for some  $k$ ,  $p'' \Vdash \llbracket h \approx k \rrbracket$  and  $p'' \Vdash \llbracket k \varepsilon b \rrbracket$ . Then for some  $p'''$  with  $p''\mathcal{R}p'''$ ,  $p''' \Vdash k \varepsilon b$ . But then  $\langle p''', k \rangle \in b$ , and so  $k = \hat{n}$  for some  $n \in \omega$ , and  $p''' \Vdash \llbracket \hat{n} \in a \rrbracket$  (definition of  $b$ ). We also have  $p'' \Vdash \llbracket h \approx \hat{n} \rrbracket$ , and it follows that  $p''' \Vdash \llbracket \neg(\hat{n} \in a) \rrbracket$ , a contradiction.

Additional changes resulting from the correction described above.

Page 264, second paragraph should begin: “Let  $C = \{a \mid a \subseteq \mathcal{G} \times \{\hat{0}, \hat{1}, \dots\}\}$ .”

Page 264, last paragraph of the Proof of Lemma 5.2, second sentence. This should begin: “Since  $a \subseteq \mathcal{G} \times \{\hat{0}, \hat{1}, \dots\}$ ...”

Page 265, paragraph following Proposition 5.5. This should read: “We are finished investigating  $\mathcal{P}(\mathcal{G} \times \{\hat{0}, \hat{1}, \dots\})$  and its subset  $C_0$ .”

**Page 272** (Problem found by Grigori Mints). The proof of Proposition 20.4.1 begins by noting that  $f \in g$  and  $\llbracket f \in g \rrbracket$  are equivalent. This is not the case, see correction to Page 228. However, the atomic case is still straightforward. For  $f, g \in \mathcal{D}_{\mathfrak{F}}^G$ ,  $p \Vdash f \in g$  if and only if  $p \Vdash_{\mathfrak{F}} f \in g$  for every  $p$ , by the definition of  $\Vdash_{\mathfrak{F}}$ . It follows that  $p \Vdash \Box\Diamond(f \in g)$  if and only if  $p \Vdash_{\mathfrak{F}} \Box\Diamond(f \in g)$  for every  $p$ .

## The following errors and typos were reported by Jonathan Farley.

**Page 23, line 17** “aet” should be “set”

**Page 25, line 14** “A2” should be “A5”

**Page 30, line 19** “qualify” should be “qualify”

**Page 38, line 2** It is true as written, but “ $y \subset x$ ” should be “ $y \subseteq x$ ”

**Page 40, line 11** “bounded” should be “non-empty bounded”

**Page 44, line 15** Do we know  $c$  is a set? Response: standard mathematical practice treats this as a set, but technically it is not justified until the Axiom of Substitution is introduced, **Ax 8** on page 170.

**Page 49, line 14** “ $x'Rb'$ ” should be “ $b'Rx'$ ”

**Page 49, line 15** “ $x \leq b$ ” should be “ $b \leq x$ ”

**Page 58, line -2** “ $M$ ” should be “ $N$ ”

**Page 59, line -5** “ $M$ ” should be “ $S$ ”

**Page 60, line -8** “ $N$ ” should be “ $\cup N$ ”

**Page 60, line -5** (Further corrected by Rolf Rolles) This should read “Every successor element of  $N$  is  $F(a)$  for some  $a \in \cup N$ .”

**Page 60, Exercise 4.3** Delete this exercise. Here is a counter-example. Take any set  $x$ , and consider  $S = \{x\}$ . The axiom of choice is not needed to say  $S$  has a choice function, but  $\cup S = x$ , and this need not have a choice function

**Page 62, line 9** The semicolon should be a comma

**Page 62, Lemma 5.4** Add the assumption that  $A \neq \emptyset$

**Page 62, line -2** “5.3” should be “5.2”

**Page 63, line 5** Exercise 4.3 has been deleted

**Page 63, line 7** After “denumerable” add “or finite”

**Page 64, line -9** Before “set of finite character” insert “non-empty”

**Page 66** It would have helped to point out at the beginning of §7 that  $g$  is defined on all sets

**Page 73, line 3** “10.2” should be “10.3”

**Page 79, following Definition 1.1** Should begin, “In general,  $F(x) \neq F''(x)$ ”

**Page 79, Second paragraph following Definition 1.1** Should contain “whereas  $F''(x)$  is”

**Page 80, line 2** Add that  $\varphi_2$  is  $1 - 1$

**Page 80, proof of Proposition 1.3, second line** “ $L$ ” should be “ $L_{<}$ ”

**Page 80, line -2** “onto” should be “into”

**Page 88, line -7** “isomporhic” should be “isomorphic”

**Page 89, Theorem 6.1** This should begin “For any functions  $f(x), g(x)$  on ordinals, and any function  $h(x, y, z)$ , where  $y$  and  $z$  are ordinals, . . .

**Page 93, O<sub>6</sub>** “Since  $x$ ” should be “Since  $S$ ”

**Page 97, last line of Proof of  $Q_4$**  “has rank  $< \alpha$ ” should be “has rank  $\leq \alpha$ ”

**Page 97, line -8** “ $\mathbb{F}$ ” should be “ $\mathbb{F}$ ”

**Page 99, next to last line of Example** “every subclass” should be “every non-empty subclass”

**Page 99, line -14** “Zermelo Fraenkel” needs a hyphen

**Page 101, line -16** “ $x$  of  $A$ ” should be “ $x$  of  $A - B$ ”

**Page 102, line -6** “set” should be “class”

**Page 306, Tarski, A. (1955)** “lattice-theoretical theorem” should be “lattice-theoretical fixpoint theorem”

## The following corrections are due to Fausto Barbero.

**Page 104** In the proof of Theorem 4.6, it should be  $y = g''x$ , not  $y = g(x)$ .

**Page 116** This is not actually a correction, but an elucidation. The second paragraph on the page, concluding  $A \cong A \times A$  from  $M \cong A$  and  $M \cong M \times M$ , does not look obvious at all to me. I give a short proof.

First, we observe that  $A \times A \cong [(A \setminus M) \cup M] \times [(A \setminus M) \cup M] = [(A \setminus M) \times (A \setminus M)] \cup [(A \setminus M) \times M] \cup [M \times (A \setminus M)] \cup [M \times M]$ . But  $M \cong A$  implies  $A \setminus M \preceq M$ . Using this injection, it is immediate to prove that each of  $[(A \setminus M) \times (A \setminus M)]$ ,  $[(A \setminus M) \times M]$  and  $[M \times (A \setminus M)]$  is  $\preceq M \times M$ . Then, using the law of additive absorption (Thm 6.2) in the first equation,  $A \times A \cong M \times M \cong M \cong A$ .

**Page 123** In the first line of the proof of Proposition 9.8, instead of  $c$  there should be  $x$ .

**Page 124** Line -3: instead of  $z = y \cup \omega$  it should be  $z = \mathcal{P}(y \cup \omega)$ .

**Page 132** Subsection “A-ranks”, third line: instead of  $p(x) \cap A$  it should be  $\mathcal{P}(x) \cap A$ .

**Page 133** Last line on page: instead of  $F_\alpha(y)$  it should be  $f_\alpha(y)$ .

**Page 139** Proof of 6.7, first line: “Let  $A \dots$ ” should be “Let  $A_0 \dots$ ”.

**Page 144** Paragraph beginning “One of the early major results. . . ,” the reading should be “. . . if  $\varphi$  is satisfiable in any INFINITE relational system at all, then it is satisfiable in a denumerable relational system” (remember that “denumerable” means having cardinality  $\omega$ ).

### Page 153

- Line 4 of (4): instead of  $w_n$ , it should be  $W$ .
- Lines 17, 19 and -7 of (4): terms of the form  $\cup \Gamma''(w_i)$  are used, but they are not well defined, since  $\Gamma$  is an  $n$ -ary function. Could one use  $\cup \Gamma''(w_i, \dots, w_i)$  instead?
- Not an error, but it is worth noting that  $n$  is used for two distinct purposes in this proof.

**Page 163** The formula (29) seems incorrect, because it should involve two (in general) distinct pairs  $\langle x_1, x_2 \rangle, \langle x_1, x_3 \rangle$ , but this is not possible if  $x_1, x_2, x_3$  are all extracted from members of one and the same  $z$  (remember that  $z$  should be one of the pairs that constitute the function  $x$ ). The correct prefix of the formula should be:

$$(\forall z_1 \in x)(\forall z_2 \in x)(\forall w_1 \in z_1)(\forall w_2 \in z_2)(\forall x_1 \in w_1)(\forall x_2 \in w_1)(\forall x_3 \in w_2)$$

**Page 167** I have the impression that the formula used to define  $L^n \cap c$  is based on the (generally wrong) assumption that  $L^n \cap c$  is transitive. It should be replaced by  $x \in c \wedge (\exists x_1 \in_2 c)(\exists x_2 \in_4 c) \dots (\exists x_n \in_{2n} c)(x = \langle x_1, \dots, x_n \rangle)$ , where  $y \in_k z$  abbreviates

$$\exists y_1 \dots \exists y_{k-1} (y \in y_1 \wedge y_1 \in y_2 \wedge \dots \wedge y_{k-1} \in z).$$

**Page 170** It is not specified anywhere that the axioms are the universal closures of the indicated formulas. Furthermore, in Axiom 2,  $\varphi(x, y_1, \dots, y_n)$  should be replaced (twice) by  $\varphi(y, y_1, \dots, y_n)$ .

**Page 176** Proof of 3.4: replace “non-empty member of  $L$ ” with “non-empty member  $a$  of  $L$ ”.

**Page 183** A further lemma needed to show that formulas stay  $\Sigma$  if the symbol  $\uparrow$  is used.

**Page 186** In clause (0) there are two occurrences of substitution of a set inside a term; but this has not (yet) been defined.

**Page 187** In the definition of  $x(y) = z$  just before Proposition 3.9, “ $\varphi(y) = z$ ” should be replaced by “ $\varphi(y) =$  the formula whose code is  $z$ ”. In “Valuations and truth”, the index  $j$  is introduced but never used.

**Page 189** Line 13, twice: replace “absolute” with “absolute over  $L$ ”. However, the same argument (and thus Theorem 4.1) should work for any first-order universe. Indeed, on **Page 190**, line -9, Theorem 4.1 is applied to a generic first-order universe  $K$ .

**Page 196** I think condition  $C_1$  should be  $K_0 = \emptyset$ . Otherwise it would be impossible that  $\pi''(K_0) = L_0$ , as claimed in Exercise 2.1. In condition  $C_3$ ,  $\alpha < \delta$  should appear as an index.

**Page 210** Line -3: The exclamation mark is unintentional.

**Page 217** Before Prop. 4.4: The last occurrence of “than” should be “them”.

**Page 219** Line 3: The argument also uses the Converse Barcan Formula.

**Page 242** Most likely, the proof of 6.1 should begin “ $u = \{q, a\}$ ” and “ $v = \{q', b\}$ ”.

**Page 256** Proof of Lemma 6.1, penultimate paragraph: “Then both...”: this is not really an error, but the existence of a common  $p'$  satisfying both the conditions described is not straightforward as the text seems to imply. I think one should proceed as follows: from the fact that  $p \Vdash \llbracket (\exists y)(y \text{ is } f(\hat{\gamma}_1) \wedge y \in s) \rrbracket$  and  $p \Vdash \llbracket (\exists y)(y \text{ is } f(\hat{\gamma}_2) \wedge y \in s) \rrbracket$  we get (by Prop. 4.3 (3)) that  $p \Vdash \llbracket (\exists y)(y \text{ is } f(\hat{\gamma}_1) \wedge y \in s) \wedge (\exists y)(y \text{ is } f(\hat{\gamma}_2) \wedge y \in s) \rrbracket$ . Then, applying prenex transformations, we obtain that  $p \Vdash \llbracket (\exists y)(\exists y')((y \text{ is } f(\hat{\gamma}_1) \wedge y \in s) \wedge (y' \text{ is } f(\hat{\gamma}_2) \wedge y' \in s)) \rrbracket$ . Applying twice the embedding property **P**<sub>11</sub>, we obtain  $p', a, b$  such that  $p' \Vdash \llbracket (a \text{ is } f(\hat{\gamma}_1) \wedge a \in s) \wedge (b \text{ is } f(\hat{\gamma}_2) \wedge b \in s) \rrbracket$ . Using once more Prop. 4.3 (3) we obtain the desired conclusion.

**Page 257** In order to account for the special case CH, in both centered formulas of the proof of 6.2 " $\hat{\omega} \in \alpha$ " should be replaced by " $\hat{\omega} \in \alpha \vee \hat{\omega} \approx \alpha$ ".

**Page 260** Two lines before the "Exercises": instead of  $\omega$  it should be  $\hat{\omega}$ .

**Page 299** Line -12: "computibility" should be "computability".