

# Notes on ‘Set Theory & Continuum Problem’

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## 1 Errata or what looks like these

p. 137. Theorem 6.1. of Chapter 10 seems to be incorrect as stated. One counterexample is: let  $A = B = 2$  and  $R = \in$ ,  $G = \ni$ . Then isomorphism  $F$  from  $(A, R)$  onto  $(B, G)$  looks like this:  $F(0) = 1$ ,  $F(1) = 0$ . Also,  $(A, R)$  clearly satisfies assumptions of the theorem and  $B$ , being an ordinal, is transitive. But  $F$  is not the Mostowski-Shepherdson map for  $(A, R)$ , since

$$0 = F(1) \neq F''(1^*) = F''(1) = F''(\{0\}) = \{1\}.$$

The theorem can be amended by putting  $(B, \in)$  instead of  $(B, G)$ . (In fact, this amendment seems to be assumed in the proof given for the theorem, so this might be classified as a misprint as well).

p.174–175. Lemma 3.2, Chapter 13, seems to use transitivity of  $K$  after all. When replacing formulas of the form  $a \in x_i$  with formulas  $(\exists y \in x_i)(y = a)$  it will not do to unfold the formula as  $(\exists y \in x_i)\forall z(y \in z \equiv a \in z)$ , for such a formula will be irrelevant (since it still contains  $a \in z$ ). One should rather understand the  $=$  in the formula along the lines of  $(\exists y \in x_i)\forall z(z \in y \equiv z \in a)$ , but the fact that  $\forall z(z \in y \equiv z \in a)$  defines  $y = a$  over  $K$  presupposes transitivity (or at least extensionality) of  $K$ . So, one can either mention transitivity of  $K$  as an assumption of the lemma, or, alternatively, one can include  $K_{a,j}^n$  into the types of distinguished subclasses mentioned in the hypothesis (1) of Theorem 3.1. This will not get into the way of the application of this theorem to  $L$  due to the presence of Exercise 3.1 of Chapter 12.

p. 185. The statement: ‘The four conditions of (0) can be collectively stated:  $t(m_1 \in m_2) = \{\ulcorner m_1 \urcorner, \ulcorner m_2 \urcorner\} \cap \omega$ ’ is incorrect, for, e.g.  $t(x_1 \in x_2) = \{1, 2\}$ , but  $\{\ulcorner x_1 \urcorner, \ulcorner x_2 \urcorner\} \cap \omega = \{\langle 0, 1 \rangle, \langle 0, 2 \rangle\} \cap \omega = \emptyset$ , since no natural number is an ordered pair of natural numbers. Therefore,  $\Sigma$ -condition (2) in the proof of Lemma 3.4 on this page must be replaced, e.g. by the following formula:

$$\begin{aligned} & \exists z(\forall x \in c(a))(\forall y \in c(a))(f(\langle x, y, 0 \rangle) = z \wedge \\ & \quad \wedge (\forall w \in z)\exists x'(x' = 0 \wedge (x = \langle x', w \rangle \vee y = \langle x', w \rangle))) \\ & \quad \wedge \exists y_1 y_2 y_3 (y_1 = \{x, y\} \wedge y_2 = \cup y_1 \wedge y_3 = \cup y_2 \\ & \quad \quad \wedge (\forall y_4 \in y_3)((\exists y_5 \in y_3)(y_5 = 0 \wedge (x = \langle y_5, y_4 \rangle \vee y = \langle y_5, y_4 \rangle))) \supset y_4 \in z) \\ & \quad ) \\ & ) \end{aligned}$$

(I was aiming at correctness rather than brevity). This formula is  $\Sigma$  due to the items (7), (9), (14), (24) from p. 159.

p. 275. The part of proof of Lemma 5.1, Chapter 20, given on this page employs the following transition<sup>1</sup>: from

$$q \Vdash [[c \approx x' \wedge c\varepsilon a]]$$

conclude to

$$\theta^{-1}(q) \Vdash [[\theta^{-1}(c) \approx \theta^{-1}(x') \wedge \theta^{-1}(c)\varepsilon a]],$$

using Theorem 3.2 and  $\mathfrak{C}$ -invariance of  $a$ . Now, this theorem warrants transition to

$$\theta^{-1}(q) \Vdash [[\theta^{-1}(c) \approx \theta^{-1}(x') \wedge \theta^{-1}(c)\varepsilon\theta^{-1}(a)],]$$

and  $\mathfrak{C}$ -invariance of  $a$  means that  $[[a \approx \theta^{-1}(a)]]$  is S4-valid in  $\mathcal{M}$ , but this does not give us our conclusion, since the book proves the substitutivity of  $\approx$  only with respect to formulas of classical language, which do not contain  $\varepsilon$ . This shows that the transition is unwarranted but not necessarily shows that it is wrong. I think that I can produce both a counterexample to the transition and a possible amendment to the proof.

*Counterexample.* Let  $\mathcal{G} = \{p, q, r, s\}$ , let  $\mathcal{R}$  be a transitive and reflexive closure of  $\{\langle p, q \rangle, \langle r, s \rangle\}$  and let  $\theta(p) = r$ ,  $\theta(r) = p$ ,  $\theta(q) = s$ , and  $\theta(s) = q$ . Then  $\theta$  is an automorphism and the group generated by  $\theta$  contains only  $\theta$  itself plus an equivalence function as an identity map. Also,  $\theta$  happens to be its own inverse. Consider then the following sets:

$$\begin{aligned} a &= \{\langle p, \hat{p} \rangle, \langle q, \hat{p} \rangle, \langle s, \hat{p} \rangle\} \\ b &= \{\langle r, \hat{p} \rangle, \langle q, \hat{p} \rangle, \langle s, \hat{p} \rangle\} \end{aligned}$$

It follows from Proposition 2.9, Chapter 20, that  $\theta(a) = b$ . It is also clear that both  $[[\hat{p}\varepsilon a]]$  and  $[[\hat{p}\varepsilon b]]$  are S4-valid in  $\mathcal{M}$ . One can also see that  $[[a \approx b]]$  is S4-valid in  $\mathcal{M}$ . Indeed, use Proposition 2.12, Chapter 17, and assume that for some  $p' \in \mathcal{G}$  we have  $p' \Vdash [[x\varepsilon a]]$ . Then, of course,  $x = \hat{p}$  and we have both  $p' \Vdash [[\hat{p} \approx \hat{p}]]$  and  $p' \Vdash [[\hat{p}\varepsilon b]]$ , and similarly for the other direction. So we have shown that  $a$  is invariant with respect to the group of automorphisms generated by  $\theta$  (since  $\hat{p}$ , the only  $\varepsilon$ -element of  $a$ , is of course invariant by Proposition 2.9). Now consider the two sets

$$\begin{aligned} c &= \mathcal{G} \times \{a\} \\ d &= \mathcal{G} \times \{b\} \end{aligned}$$

Again we have  $\theta(c) = d$  and due to the S4-validity of  $[[a \approx b]]$  we clearly have both the S4-validity of  $[[c \approx d]]$  and the invariance of  $c$  with respect to our group of automorphisms. So, to summarize, we have, e.g.

$$p \Vdash [[a \approx a \wedge a\varepsilon c]],$$

but we do not have

$$r \Vdash [[b \approx b \wedge b\varepsilon c]],$$

since  $b\varepsilon c$  never holds in  $\mathcal{M}$ . And given that  $\theta(a) = b$ ,  $\theta(p) = r$  and  $\theta = \theta^{-1}$ , this means that

$$\theta^{-1}(p) \not\Vdash [[\theta^{-1}(a) \approx \theta^{-1}(a) \wedge \theta^{-1}(a)\varepsilon c]].$$

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<sup>1</sup>I am using double square brackets instead of the involved brackets the authors use to denote the translation from modal to non-modal formulas

*Amendment.* We revise the lemma by adding the assumption that automorphisms in  $\mathfrak{F}$  are first-order definable in  $M$  and we begin the amended proof by establishing the following:

*Claim.* If  $a \in R_\alpha^{\mathcal{G}}$ , and  $\theta \in \mathfrak{C}_0$ , then  $\theta(a) \in R_\alpha^{\mathcal{G}}$ .

*Proof.* By induction on  $\alpha$ . Induction basis and the limit case are obvious; we consider successor case. Let  $\alpha = \beta + 1$  and let  $a \in R_{\beta+1}^{\mathcal{G}}$ . If  $b \in \theta(a)$ , then  $b = \langle \theta(p), \theta(c) \rangle$ , where  $p \in \mathcal{G}$  and  $c \in R_\beta^{\mathcal{G}}$ . By induction hypothesis,  $\theta(c) \in R_\beta^{\mathcal{G}}$ . Therefore,

$$\theta(a) \subseteq \mathcal{G} \times R_\beta^{\mathcal{G}}.$$

To arrive at our conclusion,  $\theta(a) \in R_{\beta+1}^{\mathcal{G}}$ , it remains to show that  $\theta(a) \in M$ ; given that  $\theta$  is first-order definable,  $a \in M$  and  $M$  is a first-order universe, this requires only an application of the corresponding version of substitution schema.  $\square$

Having the Claim established, we proceed more or less in the same way as in the book, until we get the following premises:

$$q \Vdash x' \varepsilon b \tag{1}$$

$$q \Vdash [[x \in b]] \tag{2}$$

$$q \Vdash [[\neg x \in \theta(b)]] \tag{3}$$

$$q \Vdash [[x \approx x']] \tag{4}$$

$$q \Vdash [[x' \varepsilon b]] \tag{5}$$

$$q \Vdash [[c \approx x' \wedge c \varepsilon a]] \tag{6}$$

$$x' \in R_\alpha^{\mathcal{G}} \cap \mathcal{D}_{\mathfrak{F}}^{\mathcal{G}} \tag{7}$$

$$\mathcal{M} \models_{S4} [[a \approx \vartheta(a)]] \quad (\text{for every } \vartheta \in \mathfrak{C}) \tag{8}$$

Now we reason as follows:

$$q \Vdash [[x' \in a]] \quad (\text{from (6)}) \tag{9}$$

$$\theta^{-1}(q) \Vdash [[\theta^{-1}(x') \in \theta^{-1}(a)]] \quad (\text{from (9) by T3.2}) \tag{10}$$

$$\theta^{-1}(q) \Vdash [[\theta^{-1}(x') \in a]] \quad (\text{from (8), (10) by T2.7 Ch. 17}) \tag{11}$$

Now, using (11), take some  $d \in \mathcal{D}^{\mathcal{G}}$  and  $r \in \mathcal{G}$  such that

$$\theta^{-1}(q) \mathcal{R} r \tag{12}$$

$$r \Vdash [[\theta^{-1}(x') \approx d \wedge d \varepsilon a]] \tag{13}$$

We then continue in the following way:

$$\theta^{-1}(x') \in R_\alpha^{\mathcal{G}} \cap \mathcal{D}_\beta^{\mathcal{G}} \quad (\text{from (7) by the above Claim \& T3.3}) \quad (14)$$

$$r \Vdash [[\theta^{-1}(x') \varepsilon b]] \quad (\text{from (13) and (14) by df of } b) \quad (15)$$

$$r \Vdash [[\theta^{-1}(x') \in b]] \quad (\text{from (15)}) \quad (16)$$

$$\theta(r) \Vdash [[x' \in \theta(b)]] \quad (\text{from (15) by T3.2}) \quad (17)$$

$$q\mathcal{R}\theta(r) \quad (\text{from (12)}) \quad (18)$$

$$\theta(r) \Vdash [[x \in \theta(b)]] \quad (\text{from (4), (17), and (18)}) \quad (19)$$

$$\theta(r) \Vdash [[\neg x \in \theta(b)]] \quad (\text{from (3) and (18)}) \quad (20)$$

So we got our contradiction in place.

A final erratum which I would like to call attention to is not a very harmful (in the sense that it does not render the results wrong) but still somewhat misleading and very pervasive habit of the authors to skip a transition to an  $\mathcal{R}$ -successor demanded by the contexts of the form  $p \Vdash [[a \in b]]$ . E. g. on p. 230 authors write: ‘Now,  $p \Vdash [[f \in h]]$ , so  $p' \Vdash [[f \in h]]$ . Then, for some  $a$ ,  $p' \Vdash [[a \approx f]]$  and  $p' \Vdash [[a \varepsilon h]]$ ’. In fact, as the authors note themselves on p. 228, it is  $f \in h$  that is equivalent to  $\exists a[[a \approx f \wedge a \varepsilon h]]$ , whereas  $[[f \in h]]$  is equivalent to  $[[\exists a(a \approx f \wedge a \varepsilon h)]]$ , and so  $P_{11}$  on p. 226 warrants only the following conclusion: Then, for some  $a$  and some  $p''$  such that  $p' \mathcal{R} p''$ :

$$p'' \Vdash [[a \approx f]] \wedge [[a \varepsilon h]].$$

This skipping of successor transitions, however, does not get into the way of conclusions the authors aim at, since conditions of the form  $[[X]]$  are monotonic with respect to accessibility relation. Still, it is worth noting that this skipping occurs (as far as I could notice) on the following pages of the book: 230, 232, 233, 234, 236, 237, 239, 243, 246, 247, 257, 275.

## 2 Misprints

p. 133 In the proof of Theorem 3.1, Chapter 10 while constructing an  $On$  sequence, authors demand that

$$f_{\alpha+1} = f_\alpha,$$

while the subsequent reasoning seems to suggest rather

$$f_{\alpha+1} = f_\alpha \cup g(f_\alpha),$$

where  $g(f_\alpha) = \{\langle x, g(f_\alpha''(x^*)) \rangle \mid x \in \Gamma_{\alpha+1} \setminus \Gamma_\alpha\}$ .

p.138. In the proof of Theorem 6.3, Chapter 10, the phrase ‘(by  $P_2$  of §2)’ needs to be replaced with ‘(by  $P_8$  of §2)’.

p.152 In the proof of Theorem 5.1, Chapter 11, the two occurrences of  $W$  in the phrase ‘(because any element... hence of  $W$ )’ need to be replaced with  $w_\beta$ .

p.170, footnote 1  $x = y$  is said to be abbreviation of  $\forall z(z \in x \equiv z \in y)$  whereas subsequent reasoning suggests rather  $\forall z(x \in z \equiv y \in z)$ .

p. 220. In the axiom (5) the last parenthesis is missing.

p. 235. It is necessary to remove hats from  $y_1 \dots y_k$  in the proof of Corollary 3.4, Chapter 17.

p. 277.  $\theta \in \mathcal{G}$  needs to be replaced with  $\theta \in \mathfrak{C}$ .

p. 296.  $a_G \in b_B$  needs to be replaced with  $a_G \in b_G$ .

### 3 Other notes

Exercise 2.1, Chapter 20, might actually require definability of the automorphisms in question. It is hard to supply a clear counterexample, but at least I could not see how to do without it when proving that the extension of  $\theta$  on  $\mathcal{D}^{\mathcal{G}}$  is onto. I was proving the following statement

$$f \in \mathcal{D}^{\mathcal{G}} \Rightarrow \exists g \in \mathcal{D}^{\mathcal{G}} (f = \theta(g))$$

by an induction on  $\alpha$  such that  $f \in R_{\alpha}^{\mathcal{G}}$ . Then in the successor case the induction hypothesis yields that for every  $h$  such that  $\langle p, h \rangle \in f$  there is an  $h' \in \mathcal{D}^{\mathcal{G}}$  such that  $h = \theta(h')$ . Since we know that  $f$  is an  $M$ -set, then the class of right projections of elements of  $f$  must be an  $M$ -set as well; and now, if we know that  $\theta$  is definable, we can apply first-order substitution schema to get that the class of their  $\theta$ -images is an  $M$ -set, too. Then it easily follows that the set

$$\{\langle \theta(p), \theta(h) \rangle \mid \langle p, h \rangle \in f\}$$

is an  $M$ -set which is a subset of  $\mathcal{D}^{\mathcal{G}}$ , and so it is in  $\mathcal{D}^{\mathcal{G}}$  itself.

Perhaps you had in mind some other solution to the exercise which does not require definability; but if you thought of the solution along the lines outlined above, then it might be a good idea to include the definability assumption into the formulation of the exercise.