

Exercise Set 11

Melvin Fitting

April 24, 2018

The relevance logic \mathbf{B} was characterized semantically in class. All questions here ask for you to use that semantics. To show non-validity, give a counter-model and explain why it is one. To show validity, establish that there can be no counter-model.

1. Show the validity or non-validity in \mathbf{B} of $((A \supset B) \wedge (A \supset C)) \supset (A \supset (B \wedge C))$.
2. Show the validity or non-validity in \mathbf{B} of $(P \supset Q) \supset ((P \wedge R) \supset (Q \wedge R))$.
3. This question amounts to a proof that \mathbf{B} is a relevance logic. That is, if $A \supset B$ is valid in \mathbf{B} then A and B must share a propositional variable. (It is taken directly from Graham Priest's book)
 - (a) Let \perp and \perp^* be a pair of non-normal worlds such that every propositional parameter is true at \perp and false at \perp^* . Suppose that $R\perp\perp\perp$, $R\perp^*\perp\perp^*$, and that each world accesses no other worlds. Show that every formula is true at \perp and false at \perp^* .
 - (b) Let w and w^* be a pair of non-normal worlds such that $Rw\perp w$ and $Rw^*\perp w^*$. Using part (a), show that: (i) if every parameter in A is true at w and false at w^* , the same is true of A ; (ii) if every parameter in B is false at w and true at w^* , the same is true of B .
 - (c) Use this to show that if $A \supset C$ is valid in \mathbf{B} , then A and C share a propositional parameter.